## Architectures for the FPGA Implementation of Online Kernel Methods

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### Computer Engineering Laboratory

- > Focuses on how to use parallelism to solve demanding problems
  - Novel architectures, applications and design techniques using VLSI, FPGA and parallel computing technology
- > Applications
  - Computational Finance
  - Signal Processing
  - Nanoscale Interfaces



### **Overview**

- Motivation
- > Kernel methods
  - Vector processor
  - Pipelined
  - Braided
  - Distributed
- Conclusion





How to beat other people to the money (latency)

- Low latency trading looks to trade in transient situations where market equilibrium disturbed
  - 1ms reduction in latency can translate to \$100M per year



 Latency also important: prevent blackouts due to cascading faults, turn off machine before it damages itself, etc

> Information Week: Wall Street's Quest To Process Data At The Speed Of Light



### Motivation (latency)

#### Exablaze Low-Latency Products





ExaLINK Fusion 48 SFP+ port layer 2 switch for replicating data typical 5 ns fanout, 95 ns aggregation, 110 ns layer 2 switch

Xilinx Ultrascale FPGA, QDR SRAM, ARM processor

ExaNIC X10 typical raw frame latency 60 bytes 780 ns

# What we can't do: ML with this type of latency

Source: exablaze.com



- Ability to acquire data improving (networks, storage, ADCs, sensors, computers)
  - e.g. hyperspectral satellite images, Big Data e.g. SIRCA has 3PB of historical trade data
- > Significant improvements in ML algorithms
  - Deep learning (model high-level abstractions in data) for leading image and voice recognition problems; support vector machines to avoid overfitting



## What we can't do: learning with this data rate



- To provide ML algorithms with higher throughput and lower latency we need
  - Low Energy so power doesn't become a constraint, operate off batteries (satellite and mobile)
  - **P**arallelism so we can reduce latency and increase throughput
  - Interface so we don't need to go off-chip which reduces speed and increases energy
  - **C**ustomisable so we can tailor entire design to get best efficiency

 Using FPGAs, develop improved algorithms and system implementations for ML

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### Linear Techniques

- > Linear techniques extensively studied
  - > Solution has form  $y = w^T x + b$
  - Use training data x to get maximum likelihood estimate of w or a posterior distribution of w
- > Pros
  - Sound theoretical basis
  - Computationally efficient
- Cons
  - Linear!
- > There is an equivalent dual representation

$$f(x) = \langle w, x \rangle + b = \sum o_i y_i \langle x_i, x \rangle + b$$



### e.g. Max Margin Hyperplane



### What do we do if given this problem?





> Map the problem to a **feature space** 

### Mapping to a Feature Space





Input Space

**Feature Space** 

- Choose high dimensional feature space (so easily separable)
- > BUT computing Φ is expensive!

### Kernel Trick



- > Kernel is a similarity function
  - defined by an implicit mapping  $\phi$ , (original space to feature space)

$$\kappa(x,x') = \phi(x)^T \phi(x') = \left\langle \phi(x), \phi(x') \right\rangle$$

- e.g. Linear kernel κ(x,x')=<x,x'>
- e.g. Polynomial kernel  $\kappa(x,x')=(1+\langle x,x'\rangle)^d$  for d=2:  $\phi(x) = (x_1^2, x_2^2, \sqrt{2x_1x_2})$
- e.g. Gaussian kernel (universal approximator)  $k(x, x') = \exp\left(-\frac{\|x x'\|^2}{2\sigma^2}\right)$ 
  - $\Phi(x)$  infinite in dimension!
- Modify linear ML techniques to kernel ones by replacing dot products with the kernel function (kernel trick)
  - e.g. linear discriminant analysis, logistic regression, perceptron, SOM, K-means, PCA, ICA, LMS, RLS, …
  - While we only describe prediction here, also applied to training equations

### Support Vector Machine





• The decision boundary:

$$\hat{\mathbf{w}}^T \mathbf{x} + w_0 = \sum_{i \in SV} \hat{\alpha}_i y_i(\mathbf{x}_i^T \mathbf{x}) + \mathbf{b}$$

$$\hat{y} = \operatorname{sign}\left[\sum_{i \in SV} \hat{\alpha}_i y_i (\mathbf{x}_i^T \mathbf{x}) + b\right]$$

$$K(\mathbf{x}, \mathbf{x}') = \boldsymbol{\varphi}(\mathbf{x})^T \boldsymbol{\varphi}(\mathbf{x}')$$

 $\mathbf{x} \rightarrow \boldsymbol{\varphi}(\mathbf{x})$ 

Never explicitly compute  $\Phi(x)$ , computing K(x,x') is O(m) e.g. poly kernel  $\Phi(x)$ , dimension (d+m-1)!/d!(m-1)! For d=6, m=100 this is a vector of length 1.6e9



- In Kernel-based learning algorithms, problem solving is now decoupled into:
  - A general purpose learning algorithm often linear (well-funded, robustness, ...)
  - A problem specific kernel (we focus on time series but kernels exist for text, DNA sequences, NLP)





### Examples are KLMS and KRLS

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- Traditional ML algorithms are batch based
  - Make several passes through data
  - Requires storage of the input data
  - Not all data may be available initially
  - Not suitable for massive datasets

- > Our approach: online algorithms
  - Incremental, inexpensive state update based on new data
  - Single pass through the data
  - Can be high throughput, low latency





### Kernel Online Algorithms

Two extensively studied types of online kernel methods:

- Kernel Least Mean Squares (KLMS)
  - O(N)
  - Converges slowly (steepest descent)
  - Takes a 'step' towards minimising the instantaneous error
  - e.g. KNLMS, NORMA

- Kernel recursive least squares (KRLS)
  - O(N<sup>2</sup>)
  - Converges quickly (Newton Raphson)
  - Directly calculates least squares solution based on previous training examples using Matrix Inversion Lemma (matrix-vector multiplication)
  - e.g. SW-KRLS

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### SW-KRLS Algorithm



The pseudo code of the SW-KRLS algorithm

Initialize  $K_0 = (1+c)I$  and  $K_0^{-1} = I/(1+c)$ . for n=1,2.. do Get  $\tilde{K}_n$  from  $K_{n-1}$  with Eq.(1) Calculate  $\tilde{K}_{n-1}^{-1}$  with Eq.(2) Get  $K_n$  with Eq.(3) Calculate  $K_n^{-1}$  with Eq.(4) Get the updated solution  $\alpha_n = K_n^{-1}Y_n$ end for

Computation complexity: O(N<sup>2</sup>).

$$\tilde{K}_{n} = \begin{bmatrix} K_{n-1} & k_{n}(x_{n}) \\ k_{n}(x_{n})^{T} & k_{m} + c \end{bmatrix}$$
(1)

$$\hat{\mathbf{K}}_{n}^{-1} = \begin{bmatrix} \mathbf{K}_{n-1}^{-1} (\mathbf{I} + \mathbf{b} \mathbf{b}^{T} \mathbf{K}_{n-1}^{-1T} g) & -\mathbf{K}_{n-1}^{-1} \mathbf{b} g \\ -(\mathbf{K}_{n-1}^{-1} \mathbf{b})^{T} g & g \end{bmatrix}$$
(2)

where 
$$b = k_{n-1}(x_n) \quad d = k_{nn} + c \quad g = (d - b^T K_{n-1}^{-1} b)^{-1}$$
  

$$K_n = \begin{bmatrix} k_{n-N,n-N} + c & p^T \\ p & \tilde{K}_{n-1} \end{bmatrix}$$
(3)

where  $p = [k(x_{n-N}, x_{n-N+1}), ..., k(x_{n-N}, x_{n-1})]^T$ 

$$K_n^{-1} = G - \frac{ff^T}{e}$$
where  $\tilde{K}_n^{-1} = \begin{bmatrix} e & f^T \\ f & G \end{bmatrix}$ 
(4)





Vector add C = A + B

> Microprocessor O(N) cycles

for (i = 0; i < N; i++) C[i] = A[i] + B[i]; > Vector processor O(1) cycle

VADD(C, A, B)

 Implemented as a custom KRLS vector processor using FPGA technology



### **Instruction Set**

#### ALL i , J AND L INDEXES RANGE FROM 1 TO N

Microcode (Opcode)	Function	Total Cycles
NOP(000)	No operation	1
BRANCH (0111)	BRANCH	4
VADD (0001)	Vector add	14
VSUB (0010)	Vector subtract	14
VMUL (0011)	Array multiply	10
VDIV (0100)	Vector divide	N+28
VEXP (0110)	Vector exponentiation	N+21
S2VE (1000)	Clone a vector N times	N+4
PVADD (1001)	N x Vector add	N+13
PVSUB (1010)	N x Vector subtract	N+13
PVMUL (1011)	N x Vector multiply	N+9
PVDOT (0101)	N x Vector dot product	N+9+10

# SW-KRLS and other kernel methods implemented efficiently using this simple instruction set











ALU 1 - adder, multiplier, exp and divider



### **Detailed Datapath**







### Performance Summary

### > SW-KRLS N=64

Platform	Power (W)	Latency (uS)	Energy (10^-5 J)
Our processor (DE5 5SGXEA7N)	2 (27)	1 (12.6)	1 (34)
DSP (TMS320C6678 <b>)</b>	1 (13)	355 (4476)	181 (6167)
CPU (i5-2400@3.1GHz)	1 (13)	16 (201)	8 (269)



- Microcoded vector processor for the acceleration of kernel based machine learning algorithms.
- Architecture is optimised for dot product, matrix-vector multiplication and kernel evaluation.
- > Features simplicity, programmability and compactness.

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### Obstacle to Pipelining

#### Dependency Problem







- > Finds D (dictionary which is subset of input vectors), and  $\alpha$  (weights) for function  $f(x) = \sum_{i=1}^{D} \alpha_i \kappa(x, d_i)$
- Is a stochastic gradient descent style kernel regression algorithm. Given a new input/output pair, {x<sub>n</sub>, y<sub>n</sub>}, weight update is:
- > 1. Evaluate κ between  $x_n$  and each entry of  $D_{n-1}$ , creating kernel vector, k.
- > 2. If max(k) <  $\mu_0$ , add  $x_n$  to the dictionary, producing  $D_n$
- > 3. Update the weights using:

$$\alpha_n = \alpha_{n-1} + \frac{\eta}{\varepsilon + k^T k} (y_n - k^T \alpha_{n-1}) k$$

) How can we chose κ,  $\mu_0$ , η and ε? We must do a parameter search.

### **Removing Dependencies**



- Training is usually:
   for (hyperparameters)
   for (inputs)
   learn\_model()
- Alternative is to find L independent problems
  - E.g. monitor L different things

- Our approach: run L independent problems (different parameters) in the pipeline
  - Updates ready after L subproblems
  - Less data transfer

```
for (inputs)
for (hyperparameters)
learn_model()
```

• Similar approach for multiclass classification (train C(C-1)/2 binary classifiers)

### High Throughput KNLMS









- > Area O(MN)
- Memory O(MN)
- Latency O(log<sub>2</sub>N+log<sub>2</sub>M)

	+ (11)	× (7)	/ (30)	exp (20)	< (4)
Operation	2MN + 2N	MN+ 4N+1	1	N	N-1
Latency	$\frac{\log_2 N +}{\log_2 M + 3}$	5	1	1	$\log_2 N$



- > Break feedforward/feedback path and sythesised with Vivado HLS
- > RIFFA 2.2.0 used for PCIe interface





### Performance

#### Core with input vector M=8 and dictionary size N=16 (KNLMS)

Implementation	Freq (MHz)	Time (ns)	Slowdown	
Float	314	3	1	
System	250	14	4	
Naive	97	7,829	2,462	
CPU (C)	3,600	940	296	
Pang et al (2013)	282	1,699	566	

- Energy efficient, Parallelism (pipelining), Integrated with PCIe and Customised (problem changed to remove dependencies)
- Can do online learning from 200 independent data streams at 70 Gbps (160 GFLOPS)



- > Demonstrated feasibility of a fully-pipelined regression engine
  - 200-stage pipeline achieves much higher performance than previous designs
  - 160 GFLOPS (70x speedup to CPU and 660x faster than our previous microcoded KRLS processor)
- > Also studied a fused, fixed-point floating point design details in paper
- > First such processor which can keep up with line speeds
  - Believe this is enabling technology for real-time ML applications

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NORMA

### Naive Online regularised Risk Minimization Algorithm

- > Finds D (dictionary which is subset of input vectors), and  $\alpha$  (weights) for function  $f(x) = \sum_{i=1}^{D} \alpha_i \kappa(x, d_i)$
- Minimise instantaneous risk of predictive error (R<sub>inst,λ</sub>) by taking a step in direction of gradient

$$f_{t+1} = f_t - \eta_t \partial_f R_{inst,\lambda}[f, x_{t+1}, y_{t+1}]\Big|_{f=f_t}$$

> Can be used for classification, regression, novelty detection

> Update for novelty detection

$$(\alpha_i, \alpha_t, \rho) = \begin{cases} (\Omega \alpha_i, 0, \rho + \eta \nu) \text{ if } f(x_t) \geq \rho & \text{Add } x_{t+1} \text{ to dictionary} \\ (\Omega \alpha_i, \eta, \rho - \eta(1 - \nu)) & \text{otherwise} \end{cases}$$

### Datapath for NORMA







### NORMA Update (Case 1)

$$\alpha_{i}, \alpha_{t}, \rho) = \begin{cases} (\Omega \alpha_{i}, 0, \rho + \eta \nu) \text{ if } f(x_{t}) \geq \rho & \text{Add } \mathbf{x}_{t+1} \text{ to dictionary} \\ (\Omega \alpha_{i}, \eta, \rho - \eta(1 - \nu)) \text{ otherwise} \end{cases}$$
New Example
$$\overbrace{\mathbf{x}_{t+1}} & \overbrace{\mathbf{x}_{t+1}} & \overbrace{\mathbf{x}_{t+$$



### NORMA Update (Case 2)

$$(\alpha_{i}, \alpha_{t}, \rho) = \begin{cases} (\Omega \alpha_{i}, 0, \rho + \eta \nu) \text{ if } f(x_{t}) \geq \rho & \text{Add } x_{t+1} \text{ to dictionary} \\ (\Omega \alpha_{i}, \eta, \rho - \eta(1 - \nu)) \text{ otherwise} \end{cases}$$
New Example
$$(d_{1} + \kappa(\cdot, \cdot) + \kappa(\cdot,$$



### Properties of NORMA

- > NORMA is a sliding window algorithm
  - If new dictionary entry added  $[d_1, \cdots d_D] \rightarrow [x_t, d_1, \cdots d_{D-1}]$
  - Weight update is just a decay  $\alpha_i \rightarrow \Omega \alpha_i$
  - Update cost is small compared to computing f(x<sub>t</sub>)
- Is this really true?







- > Recall carry select adder
  - implement both cases in parallel and select output



### Braiding



$$f(x_{t+1}) = \sum_{i=1}^{D} \alpha_i \kappa(x_{t+1}, d_i)$$

Use the previous dictionary for  $x_t$  denoted  $\hat{d}_i$ 

$$f(x_{t+1}) = \sum_{i=1}^{D-1} \Omega \hat{\alpha}_i \kappa(x_{t+1}, \hat{d}_i) + \text{something}$$

if  $x_t$  is added then this term  $= \alpha_{x_t} \kappa(x_{t+1}, x_t)$ if  $x_t$  is not added then this term  $= \Omega \alpha_D \hat{\kappa}(x_{t+1}, \hat{d_D})$ 



### **Braiding Datapath**





### Generalised to p cycles

$$\begin{split} f_t(x_{t+1}) &= \sum_{i=1}^{D-p} \Omega^p \hat{\alpha}_i \kappa(x_{t+1}, \hat{d}_i) \\ & = \begin{cases} + \begin{cases} 0 \text{ if } x_{t+1-p} \text{ is not added} \\ \Omega^{p-1} \alpha_{x_{t+1-p}} \kappa(x_{t+1}, x_{t+1-p}) \text{ otherwise} \\ + \begin{cases} 0 \text{ if } x_{t+2-p} \text{ is not added} \\ \Omega^{p-2} \alpha_{x_{t+2-p}} \kappa(x_{t+1}, x_{t+2-p}) \text{ otherwise} \\ \vdots \\ + \begin{cases} 0 \text{ if } x_t \text{ is not added} \\ \alpha_{x_t} \kappa(x_{t+1}, x_t) \text{ otherwise} \end{cases} \\ & + \sum_{i=D-p+1}^{D-q} \Omega^p \hat{\alpha}_i \kappa(x_{t+1}, \hat{d}_i) \end{split}$$

$$m_{i} = \kappa(d_{i}, x_{j})$$
(k cycles)
$$f_{t}(x_{t+1}) =$$

$$\sum_{i=1}^{D} \alpha_{i}m_{i} \text{ (s cycles)}$$

$$\alpha(f_{t}(x_{t+1}))$$
(1 cycle)

Pipeline (p cycles)





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- Implemented in Chisel
- > On XC7VX485T- 2FFG1761C achieves ~133 MHz
- > Area O(FDB<sup>2</sup>) (F=dimensionality of input vector), time complexity O(FD)
- Speedup 500x compared with single core CPU i7-4510U (8.10 fixed)

F=8, D=	16	32	64	128	200
Frequency (MHz)	133	138	137	131	127
DSPs (/2,800)	309	514	911	1,679	2,556
Slices (/759,000)	4615	8194	14,663	29,113	46,443
Latency (cycles)	10	11	12	12	13
Speedup (×)	47	91	178	344	509
Latency reduction (×)	4.69	8.30	14.9	28.7	39.2



### Comparison of Architectures

Core with input vector F=8 and dictionary size D=16

Design	Precision	Freq MHz	Latency Cycles	T.put Cycles	Latency nS	T.put nS
Vector KNLMS	Single	282	479	479	1,699	1,699
Pipelined KNLMS	Single	314	207	1	659	3.2
Braided NORMA	8.10	113	10	1	89	8.8

Open source (GPLv2): github.com/da-steve101/chisel-pipelined-olk





- Braiding: rearrangement of a sliding window algorithm for hardware implementations
  - NORMA used but other ML algorithms possible
- Compared with pipelined KNLMS,
  - 20x lower latency at 1/3 of the throughput

> Open source (GPLv2): github.com/da-steve101/chisel-pipelined-olk

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#### **Distributed KRLS**

- > One problem with KRLS is how to get scalable parallelism
- Proposed a method, which uses KRLS (Engel et al. 2004) to create models on subsets of the data.
- These models can then be combined using KRLS again to create a single accurate model
  - > We have shown an upper bound on the error introduced





### Accuracy

#### Distributed KRLS Vs Cascade SVM

Accuracy comparison







#### Distributed KRLS Vs Cascade SVM

> Average Speedup about 20x on a 16 node cluster



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### Conclusion

Demonstrated high-performance applications in ML



- > Machines of the future will need to interpret and process data using ML
  - FPGAs are a key enabling technology for energy-efficient, fast implementations
  - A lot more to do!



- Stephen Tridgell, Duncan J.M. Moss, Nicholas J. Fraser, and Philip H.W. Leong. Braiding: a scheme for resolving hazards in NORMA. In *Proc. International Conference on Field Programmable Technology (FPT)*, page to appear, 2015.
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### Thank you!



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### Australian Institute for Nanoscale Science and Technology (AINST)

#### > Type A, B and C Laboratories

- Temperature: ± 0.1 Degree (Type A) to ± 0.5 Degree (Type C)
- Humidity: ± 5% (Type A) to ± 10% (Type C)
- > Vibration: VCE (Type A) to VCB (Type C)
- > EMI: 0.3mGp-p (Type A) to 3mGp-p (Type C)







### "Pure and Bright" Photon Source

- Photons a vital resource for the implementation of quantum computing and key distribution
  - Key enabling block are photon sources
- Our approach: generate correlated photon pairs where one "heralds" existence of partner
  - excellent spatial-temporal-spectral properties but  $\mu << 1$
  - for *n* probability of successful photon-photon interaction scales as  $\mu^n$  and becomes impractically small (record is n=8)
- > This work: increase  $\mu$



## Temporal multiplexing





- 1. Precisely synchronizing photons clocks
- 2. Managing their arrival time to the accuracy of several picoseconds
- 3. Controlling their polarization
- 4. Ultra-low loss components so this is achieved while maintaining photons' indistinguishability (in frequency, temporal and polarization degrees of freedom)

Obstacles since temporal multiplexing proposed by Mower (2011)







### Are our photons single? g<sup>2</sup>(0) measurement



g<sup>2</sup>(0) is a measure of level of multi-photon noise

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### Are our photons indistinguishable? Hong-Ou-Mandel interference







### Convergence

#### RICHARD et al.: ONLINE PREDICTION OF TIME SERIES DATA WITH KERNELS



Fig. 2. Learning curves for KAP, KNLMS, SSP, NORMA and KRLS obtained by averaging over 200 experiments.